SYNTAX



Matt Post IntroHLT class 5 September 2024



So gorgeous was the spectacle on the May morning of 1910 when nine kings rode in the funeral of Edward VII of England that the crowd, waiting in hushed and blackclad awe, could not keep back gasps of admiration.

Barbara W. Tuchman, The Guns of August

morning keep could awe, the crowd, admiration. in hushed and black-clad of funeral May gorgeous of not on of rode waiting the VII England 1910 back that spectacle the Edward the in gasps kings was when nine of So

Other examples

<u>GOOD</u>		<u>BAD</u>
for i in range(args.N): print(i)	python program	i in: i for range(print)
<html> Lorem ipsum </html>	HTML	ipsum <h> Lorem <ptml></ptml></h>
http://google.com	URLs	gsd@ht//:ww 4

What are the abstractions and tools that underlie all of these examples?

Today we will cover



Goals for today

- After today, you should be able to
 - describe syntax both mathematically and linguistically
 - enumerate the formal language (Chomsky) hierarchy
 - provide a description of constituent grammars
 - sketch the algorithm for CKY parsing

Outline

formal language theory

natural language

parsing

Formal Language Theory

- Generalization: define a language to be a set of strings under some alphabet, Σ
 - e.g., the set of valid English sentences (where the "alphabet" is English words), or the set of valid Python programs

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Formal Language Theory

- Generalization: define a language to be a set of strings under some alphabet, Σ
 - e.g., the set of valid English sentences (where the "alphabet" is English words), or the set of valid Python programs
- Formal Language Theory provides a common framework for studying properties of these languages, e.g.,
 - Is this file a valid C++ program? A valid Czech sentence?
 - What is the structure? <=> How do I find its meaning?
 - How hard / time-consuming is it to answer these questions?

Languages as sets

- $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- What do you think these languages describe (in words?)

$$\mathscr{L}_1 = \{0, 1, 2, 3, 4, 5, \ldots\}$$

 $\mathscr{L}_2 = \{-12.4, 0, 142, 142.1, 142.01, 142.001, \ldots\}$

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Definitions

formal name	think	description	repr
letter	token	the fundamental unit under consideration (e.g., a word, or a UTF-8-encoded letter)	a, b,
alphabet	vocabulary	A set of tokens	Σ
word	"string"	a sequence of zero or more tokens in the vocabulary	α, β, \dots
language	language	a set of strings	Ľ

Some notes

- Σ^* is the set of all strings in a vocabulary, Σ
- One special string is the empty string, {} or ϵ
- A language \mathscr{L} can be very large—even infinite!
 - In fact, most languages probably are
 - List a few

Generative descriptions of lang.

- A definition of languages as **sets** is not very useful
 - Why not?
- A better approach:
 - Develop a process that can describe how strings in a language
 - New membership criteria:
 - IN: can be generated by this process
 - **OUT**: cannot be generated by this process

Generative examples

- $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- You probably know one generative process already: regular expressions
- What do these languages describe (in words?)

$$\mathcal{L}_1 = \Sigma^*$$

 $\mathcal{L}_2 = 0 \mid [1 - 9][0 - 9]^*$

Regular language examples

- How can we write the following languages?
 - All floating point numbers
 - Email addresses
- These are much more compact representations compared to their set notations!

Regular languages with rules

- The "regular expression" syntax is a shortcut representation
- We can describe the generative process more formally using a set of *rules* that are recursively applied

 $A \rightarrow Aa$ $A \rightarrow a$

- Rules have two types of symbols:
 - *terminal* symbols (lowercase letter) are normal vocabulary items
 - nonterminal symbols (capital letters) are recursively replaced until there are no more of them

• $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$

- $\Sigma = \{0,1,2,3,4,5,6,7,8,9\}$ • $\mathscr{L}_1 = \Sigma^*$ $S \rightarrow A$
 - $A \rightarrow 0A$
 - $A \rightarrow 1A$
 - $A \rightarrow 2A$

$$A \rightarrow 9A$$

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- $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$ • $\mathscr{L}_1 = \Sigma^*$ S \rightarrow A A \rightarrow 0A A \rightarrow 1A A \rightarrow 2A
 - ••••
 - $A \rightarrow 9A$
- $\cdot A \rightarrow \epsilon$

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- $\mathcal{L}_2 = 0 \mid [1 9][0 9] *$

- $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- $\mathcal{L}_2 = 0 \mid [1 9][0 9] *$
- · S→A
 - A→0
 - A→1B
 - A→2B
 - • •
 - A→9B
 - B→0B
 - B→1B
 - ...
 - $B \rightarrow 9B$ $B \rightarrow \epsilon$

Formal definition of a language

- Definitions: consider the set $(\Sigma, N, S \in N, R)$, where
 - Σ is the *vocabulary* which is a finite set of *terminal symbols*
 - N is a finite set of *nonterminals symbols*
 - $S \in N$ is a special nonterminal called the *start symbol*
 - $\alpha, \beta, and \gamma$ are *strings* of zero or more terminal and nonterminal symbols
 - R is a set of rules of the form $\alpha N\beta \rightarrow \gamma$

Regular languages

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 ightarrow \gamma$

Туре	Rules	Name	Recognized by
3	A → aB	Regular	Regular expressions

 All the languages we created earlier (for example, the set of email addresses) can be described with such rules

Context-free languages

- Definitions: consider the set $(\Sigma, N, S \in N, R)$, where
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 - α , β , and γ are strings of zero or more terminal and nonterminal symbols
 - R is a set of rules of the form $\alpha N\beta \to \gamma$

Туре	Rules	Name	Recognized by
2	$A \rightarrow \alpha$	Context-free	Pushdown automata

 This change might seem small, but it fundamentally alters the kinds of languages that can be generated

Context-free and not regular

•
$$\Sigma = \{a, b, c, \dots, z\}$$

- Create a context-free language for ${\mathscr L},$ the set of palindromes

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- · S→A
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 - $A \rightarrow z A z$ $A \rightarrow \epsilon$

Context-free and not regular

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$$\Sigma = \{a, b, c, \dots, z\}$$

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- · S→A
 - A→aAa
 - A→bAb

 $A \rightarrow z A z$ $A \rightarrow \epsilon$

 Can you do this with the regular language constraint on rules?



The Chomsky Hierarchy

- Named after Noam Chomsky, the MIT linguist
- Different constraints on the rules lead to more powerful sets of languages that can be described
- More powerful languages are harder (meaning, more compute-intensive) to recognize

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Туре	Rules	Name	Recognized by	Complexity
3	A → aB	Regular	Regular expressions	$\mathcal{O}(n)$
2	$A \rightarrow \alpha$	Context-free	Pushdown automata	$\mathcal{O}(n^3)$
1	$\alpha A \beta \rightarrow \alpha \gamma \beta$	Context- sensitive	Linear-bounded Turing machine	$\mathcal{O}(2^n)$
0	$\alpha A \beta \rightarrow \gamma$	Recursively enumerable	Turing Machines	undecidable

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- Consider, for example, *computer programs*
 - They either compile or don't compile
 - Their structure determines their interpretation

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- We can generalize this discussion to make a connection between natural and other kinds of languages
- Consider, for example, *computer programs*
 - They either compile or don't compile
 - Their structure determines their interpretation
- What is the structure?

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Linguistic fields of study

Phonetics: sounds
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- Phonology: sound systems

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- Semantics: sentence meaning
- Pragmatics: contextualized meaning and communicative goals

Today's focus

BENDER

LINGUISTIC FUNDAMENTALS FOR NATURAL LANGUAGE PROCESSING

yntax

-who did what to rrse of the problem lationship between be extremely useful to NLP. Likewise, systems and other laccessible fashion s that can be useful us more successful

78-1-62705-011

MORGAN & CLAYPOOL PUBLISHERS

Linguistic Fundamenta for Natural Language Processing

1100 Esssemttraalls fromm

Morphology and Symtax

Emily M. Bender

- Excellent book
- Organized into 100 minilectures
- PDF available for free via JHU library (along with tens of others in the series)
- <u>https://tinyurl.com/</u> <u>linguistic-fundamentals</u>

MORGAN&CLAYPOOL

Synthesis Lectures on Human Language Technologies

Graeme Hirst, Series Editor

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What is syntax?

- A set of constraints on the possible sentences in the language
 - *A set of constraint on the possible sentence.
 - *Dipanjan had [a] question.
 - *You are <u>on</u> class.
- At a coarse level, we can divide all possible sequences of words into two groups: *valid* and *invalid* (or *grammatical* and *ungrammatical*)

- No general agreement about the exact set of parts of speech
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 - nouns: NN, NNS, NNP, NNPS
 - adverbs: RB, RBR, RBS, RP
 - verbs: VB, VBD, VBG, VBN, VBP, VBZ
 - (Here, different tags are used to capture the small bit of morphology present in English)

Parts of Speech (POS)

 \cdot Three definitions of \mathbf{noun}

<u>Grammar school</u> ("metaphysical") *a person, place, thing, or idea*

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Distributional

the set of words that have the same distribution as other nouns

{*l,you,he*} saw the {*bird,cat,dog*}.

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Functional

the set of words that serve as arguments to verbs



Phrases and Constituents

- Longer sequences of words can perform the same function as individual parts of speech:
 - I saw [a_{DT} kid_N]_{NP}
 - I saw [a kid playing basketball]_{NP}
 - I saw [a kid playing basketball alone on the court]NP
- This gives rise to the idea of a *phrasal constituent*, which functions as a unit in relation to the rest of the sentence

Constituent tests

- How do you know if a phrase functions as a constituent?
- A few tests
 - Coordination
 - Kim [read a book], [gave it to Sandy], and [left].
 - Substitution with a word
 - Kim read [a very interesting book about grammar].
 - Kim read [it].
 - See Bender #51

Constituent structure

- The head often constrains the internal structure of a constituent
- Examples
 - verb
 - [Kim]^{ARGUMENT} is [ready]^{ADJUNCT}.
 - adjective
 - Kim is [ready_{ADJ} [to make a pizza]_V].
 - Kim is [tired_{ADJ} [to make a pizza]_V].
 - noun
 - [The [red]_{ADJ} ball]
 - * [The [red]_{ADJ} ball [the stick]_N]
 - [The [red]_{ADJ} ball [on top of the stick]_{PP}]

More examples

- Kim planned [to give Sandy books].
- * Kim planned [to give Sandy].
- Kim **planned** [to give books].
- * Kim planned [to see Sandy books].
- Kim [would [give Sandy books]].
- Pat [helped [Kim give Sandy books]].
- * [[Give Sandy books] [surprised Kim]].

Human judgments

- How do we know what's in and out? We simply ask humans
- But how do humans know? This is the tie-in to formal language theory

 A finite set of rules licensing a (possibly infinite) number of strings Chomsky formal language hierarchy refresher

finite state machine
context-free grammar
context-sensitive grammar
Turing machine

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 - [noun phrase] → [determiner]?
 [adjective]* [noun]

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 [adjective]* [noun]
 - [predicate] → [verb phrase] [adjunct]

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- Rules are *phrasal* or *terminal*

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 - [subject] → [noun phrase]
 - [noun phrase] → [determiner]?
 [adjective]* [noun]
 - [predicate] → [verb phrase] [adjunct]
- Rules are *phrasal* or *terminal*
 - Phrasal rules form **constituents** in a tree
 - Terminal rules are parts of speech and produce words

Chomsky formal language hierarchy refresher



Example	
S → NP VP .	S → NP VP

 $S \rightarrow NP VP .$ $S \rightarrow [JJ NNS] VP .$

$S \rightarrow NP VP$ $VP \rightarrow JJ NNS$

 $S \rightarrow NP VP$. $S \rightarrow [JJ NNS] VP$. $S \rightarrow [Human] NNS VP$. $S \rightarrow NP VP$ $VP \rightarrow JJ NNS$ $JJ \rightarrow Human$

 $S \rightarrow NP VP$.

 $S \rightarrow [JJ NNS] VP$.

 $S \rightarrow [Human] NNS VP$.

 $S \rightarrow$ Human [languages] VP .

 $S \rightarrow NP VP$ VP $\rightarrow JJ NNS$ JJ \rightarrow Human NNS \rightarrow languages

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- $S \rightarrow NP VP$.
- $S \rightarrow [JJ NNS] VP$.
- $S \rightarrow$ [Human] NNS VP .
- $S \rightarrow$ Human [languages] VP .
- $S \rightarrow$ Human languages [VBP ADJP] .

 $S \rightarrow NP VP$ $VP \rightarrow JJ NNS$ $JJ \rightarrow Human$ $NNS \rightarrow languages$ $VP \rightarrow VBP ADJP$

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- $S \rightarrow$ Human languages are [hard] SBAR .
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- $S \rightarrow$ Human languages are [hard] SBAR .
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- $S \rightarrow$ Human languages are hard [TO VP] .
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- $S \rightarrow$ Human languages [are] ADJP .
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- $S \rightarrow$ Human languages are [hard] SBAR .
- $\mathsf{S} \to \mathsf{Human}$ languages are hard [VP] .
- $\mathsf{S} \to \mathsf{Human}$ languages are hard [TO VP] .
- $S \rightarrow$ Human languages are hard [to] VP .
- $\mathsf{S} \to \mathsf{Human}$ languages are hard to [VB] .
- $S \rightarrow$ Human languages are hard to [parse].
- $S \rightarrow$ Human languages are hard to parse .



Treebanks

- Collections of natural text that are annotated according to a particular syntactic theory
 - Usually created by linguistic experts
 - Ideally as large as possible
 - Theories are usually coarsely divided into *constituent/ phrase* or *dependency* structure

Penn Treebank (1993)

ABOUT MEMBERS

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LR Wiki	
DATA MANAGEMENT	
COLLABORATIONS	

Home > Language Resources > Data

Treebank-3

Item Name:	Treebank-3			
Author(s):	Mitchell P. Marcus, Beatrice Santorini, Mary Ann Marcinkiewicz, Ann Taylor			
LDC Catalog No.:	LDC99T42			
ISBN:	1-58563-163-9			
ISLRN:	141-282-691-413-2			
Member Year(s):	1999			
DCMI Type(s):	Text			
Data Source(s):	telephone speech, newswire, microphone speech, transcribed speech, varied			
Project(s):	TIDES, GALE			
Application(s):	parsing, natural language processing, tagging			
Language(s):	English			
Language ID(s):	eng			
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Citation:	Marcus, Mitchell, et al. Treebank-3 LDC99T42. Web Download. Philadelphia: Linguistic Data Consortium, 1999.			
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Introduction

This release contains the following Treebank-2 Material:

- One million words of 1989 Wall Street Journal material annotated in Treebank II style.
- A small sample of ATIS-3 material annotated in Treebank II style.
- A fully tagged version of the Brown Corpus.

and the following new material:

- · Switchboard tagged, dysfluency-annotated, and parsed text
- Brown parsed text

The Treebank bracketing style is designed to allow the extraction of simple predicate/argument structure. Over one million words of text are provided with this bracketing applied.

Data

The Penn Treebank

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- Contains 74 total tags: 36 parts of speech, 7 punctuation tags, and 31 phrasal constituent tags, plus some relation markings
- Was the foundation for an entire field of research and applications for over twenty years

((S

```
(NP-SBJ
 (NP (NNP Pierre) (NNP Vinken))
 (, ,)
 (ADJP
  (NP (CD 61) (NNS years))
  (JJ old))
 (, ,))
(VP (MD will)
 (VP (VB join)
  (NP (DT the) (NN board))
  (PP-CLR (IN as)
   (NP (DT a) (JJ nonexecutive) (NN director) ))
  (NP-TMP (NNP Nov.) (CD 29) )))
(...)))
```

Pierre Vinken, 61 years old, will join the board as a nonexecutive director Nov. 29.



Pierre Vinken, 61 years old, will join the board as a nonexecutive director Nov. 29.

Summary

- Formal language theory is a theory that does the following:
 - provides a compact representation of a language
 - provides an account for how strings within a language are generated
- It's very useful for describing many simple languages
- It can also be applied to natural language

Outline

formal language theory

natural language

parsing



Where we are

- We discussed formal language theory
- We showed how it might apply to human language
- But how do we get a computer to use it?
 - Sentences (or other strings we wish to parse) are observed; the structure is hidden
 - We assume these were generated by a model
 - We need
 - An algorithm for finding the sequence of actions under that model, most likely to have produced it
 - A way to learn that model

Where do grammars come from?

Where do grammars come from?

- Treebanks!
- Given a treebank, and a formalism, we can learn statistics by counting over the annotated instances



I stole this joke from Chris Callison-Burch

https://www.shutterstock.com/image-vector/stork-carrying-baby-boy-133823486

Probabilities

- For example, a context-free grammar
- \cdot We can get probabilities by reading all instances from a Treebank

$$P(A \rightarrow B \ C) = \sum_{A' \in N} \frac{P(A)}{P(A')} \quad \leftarrow a \ CFG \ rule \\ \leftarrow all \ CFG \ rules \ with \ the \ same \ left hand \ side$$

• e.g.,

- S \rightarrow NP , NP VP .	[0.002]
- NP → NNP NNP	[0.037]
$-$, \rightarrow ,	[0.999]
- NP → *	[X]
- VP → VB NP	[0.057]
- NP → PRP\$ NN	[0.008]
→.	[0.987]

Parsing

- If the grammar has certain properties (Type 2 or 3), we can efficiently answer the first question (find the hidden structure) with a **parser**
 - **Q1**: is the sentence in the language of the parser?
 - **Q2**: What is the structure above that sentence?

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Algorithms

The CKY algorithm for parsing with constituency grammars

Chart parsing for constituency grammars

- Maintains a chart of nonterminals spanning words, e.g.,
 - NP over words 1..4 and 2..5
 - VP over words 4..6 and 4..8
 - etc
- Build this chart from the bottom upward: the *opposite* direction from generation



- How do we produce this chart? Cocke-Younger-Kasami (CYK/ CKY)
- Basic idea is to apply rules in a bottom-up fashion, applying all rules, and (recursively) building larger constituents from smaller ones
- Input: sentence of length N for width in 2..N

```
for begin i in 1..{N - width}
```

```
j = i + width
```

```
for split k in \{i + 1\}..\{j - 1\}
```

```
for all rules A \rightarrow B C
```

```
create iAj if iBk and kCj
```













CKY a	Igorithr	η			
		$VP \rightarrow VB$	PP		
NP→	$\frac{VP \rightarrow _2 VB_3 \ _3 NP_5}{PP \rightarrow _2 IN_3 \ _3 NP_5}$				
$NP \rightarrow NN$	1		NP-	DT NN	
NN	NN, VB	VB,IN	DT	NN	
Tim 0	e flies	like 2 3	an	arrow	5 ₅₄



CKY algorithm						
	$S \rightarrow _{0}NP_{2} _{2}VP_{5}$ $S \rightarrow _{0}NP_{1} _{1}VP_{5}$					
	VP→VB PP					
NP→NN NN	$VP \rightarrow _2 VB_3 _3 NF$ $PP \rightarrow _2 IN_3 _3 NP$	D 5				
$NP \rightarrow NN$	NP→D	TNN				
NN NN,VB	VB,IN DT	NN				
Time flies	like an 2 3 4	arrow 5 54				

- Parsing questions:
 - **Q1**: is a given sentence in the language of the parser?
 - **Q2**: What is the structure above that sentence?
- Termination: is there a chart entry at $_0S_N$?
 - ✓ string is in the language (Q1)
 - Structures can be obtained by following backpointers in dynamic programming chart (not covered today)
- Other technical details not covered today:
 - The probability of each parse is the product of the rule probabilities
 - Ambiguities are resolved with these scores
- Demos:
 - Berkeley Neural Parser: <u>https://parser.kitaev.io</u>
 - Spacy dependency parser: <u>https://explosion.ai/demos/</u> <u>displacy</u>

Summary

